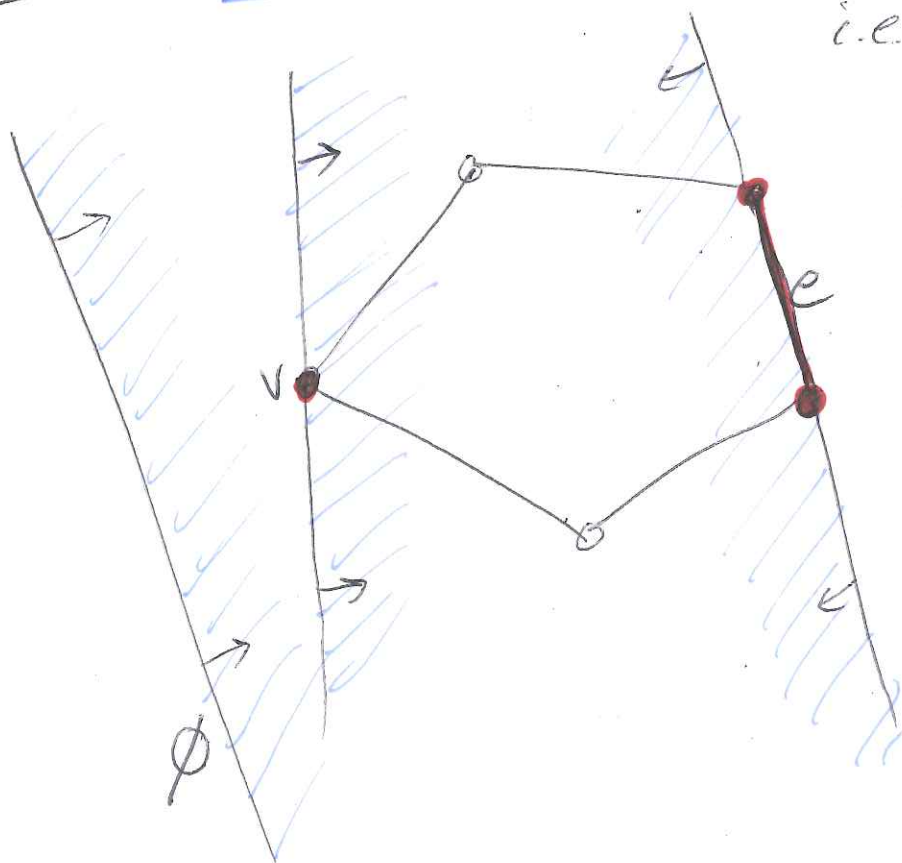


Def For a convex set  $S$ , a supporting hyperplane  $H$  is one that has all of  $S$  in one closed halfspace bounded by  $H$ .

Def A face of a polyhedron  $P$  is the intersection of  $P$  with a supporting hyperplane. We also say (by convention) that  $P$  is a face of itself.

Def A facet is a face of codimension 1.  
i.e.  $\text{dimension} = d - 1$ .



Note: The polyhedral complex gives a poset

Def For a polyhedral complex  $K$ ,  
the underlying space  $|K| = \bigcup_{P \in K} P$ .

Def We'll say  $K$  is a decomposition of  $S$   
if  $S = |K|$ .

Convex

# Decomposition from a point set

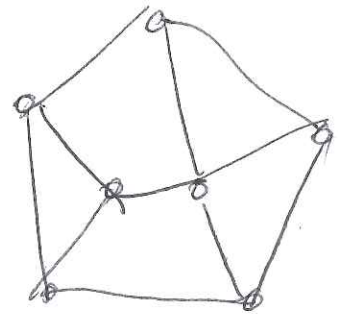
Input:  $P \subset \mathbb{R}^2$

Output: Polyhedral complex  $K$  s.t.

$P = 0\text{-faces of } K$  and

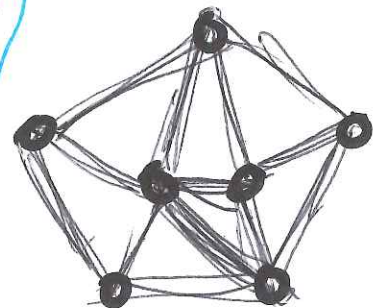
$$|K| = \text{CC}(P).$$

Convex closure



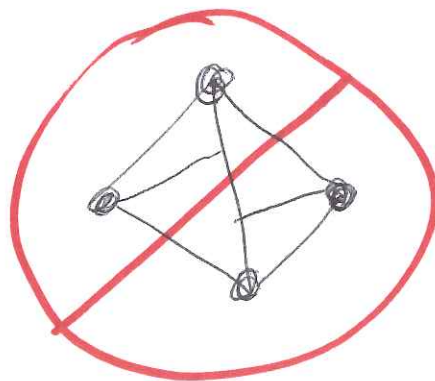
We say a convex decomposition of  $P$  is a triangulation of  $P$  if every 2-face is a triangle.

Sometimes denoted  $\Delta^n$



Note: for  $P \subset \mathbb{R}^d$ , we say  $K$  is a  $\Delta^n$  if every  $d$ -face is a simplex.

convex closure of  $d+1$  affinely independent points

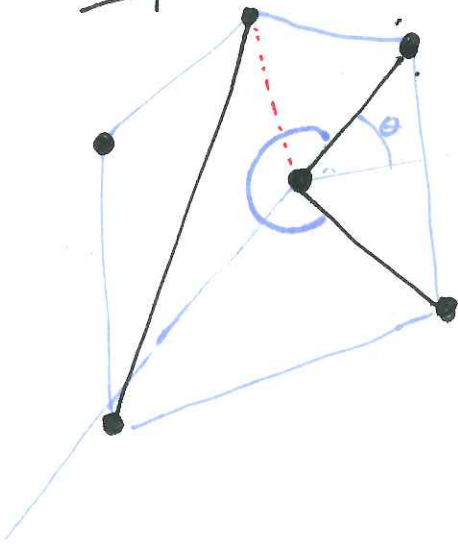


# Triangulations Exist

A Paper and Pen algorithm for triangulating a point set:

If you can draw a straight edge between two vertices that doesn't cross any other edges, then draw it. Repeat.

Why does this give a  $\Delta^n$ ?



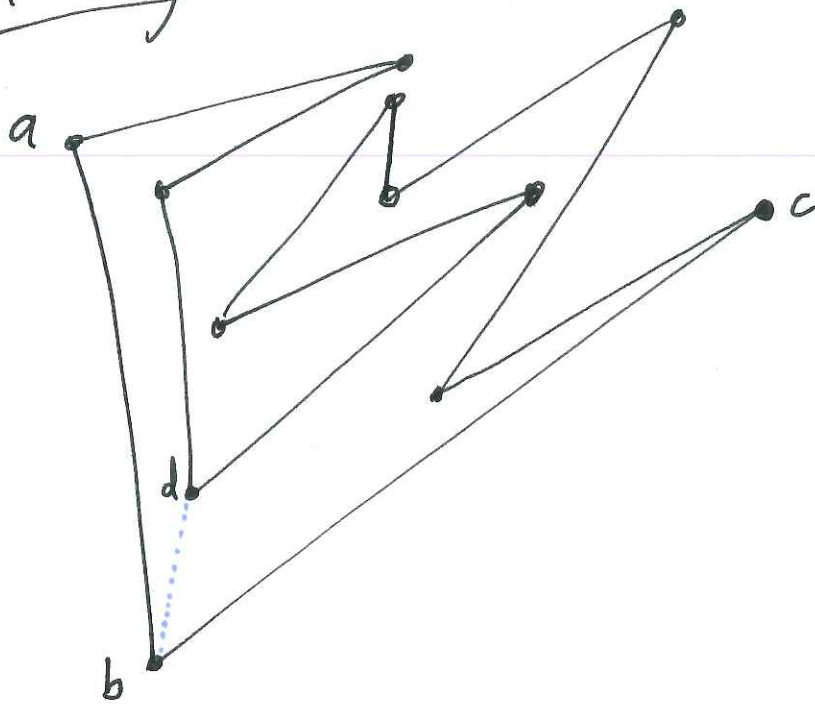
First, observe that the C.H. is in the output.

Next, for each interior vertex and angle  $\theta$  there must be an edge in the range of directions  $[\theta, \theta + \pi]$ .

So, all faces are convex polygons.

Last, convex polygons can be triangulated.

# Triangulations of Simple Polygons Exist



Idea: pick any non-reflex angle  $abc$ .  
If  $\triangle abc$  is empty, add edge  $\overline{ac}$ .  
If  $\triangle abc$  non empty, add edge  $\overline{ad}$  where  $d$  is the vertex in  $\triangle abc$  closest to edge  $\overline{ab}$ .