Homework 1

- 1. Did you do the reading? YES/NO/SORTA
- 2. Did you do the reading before class? YES/NO/SORTA
- 3. How long did you spend on this homework (rounding up)? \_\_\_\_\_hours.

## 1 Sets

**Objective:** Read and write formal descriptions of sets. Manipulate sets using basic set operations. Give a formal description of the following sets.

- The set of prime numbers less than 15.
- The set consisting of the strings *aba* and *baa*.
- The set of integers less than 8.
- The set containing the empty string.
- The set that doesn't contain anything at all.

Let  $X = \{1, 2, 3, 4\}$  and let  $Y = \{2, 4\}$ .

- Is X a subset of Y?
- Is Y a subset of X?
- What is  $X \cap Y$ ?
- What is  $X \cup Y$ ?
- What is  $Y \times X$ ?
- What is  $X \setminus Y$ ?
- What is the power set of Y?

## 2 Proofs

**Objective:** Write clear, correct proofs using different techniques. Identify flawed arguments.

The questions in this section all deal with something called graph isomorphism. Let G = (V, E) and H = (U, D) be two undirected graphs. An isomorphism is a bijection  $f : V \to U$  such that  $\{u, v\} \in E$  if and only if  $\{f(u), f(v)\} \in D$  (Recall that a bijection is a one-to-one, onto function).

If there exists an isomorphism between the vertex sets of graphs G and H, we say that G and H are *isomorphic* and denote this  $G \cong H$ .

• Prove that  $\cong$  is an equivalence relation. Recall, this means that  $\cong$  is symmetric, reflexive, and transitive.

• Find the bug in the following inductive proof that in any set of h graphs, every pair is isomorphic. Express your answer in one clear, concise sentence.

**Base case:** If h = 1. In any set containing just one graph, the one graph is isomorphic to itself since  $\cong$  is reflexive.

**Inductive Step:** For  $k \ge 1$ , assume that the claim is true for h = k and prove that it is true for h = k + 1. Take any set X of k + 1 graphs. We show that every pair of graphs in X is isomorphic. It will suffice to show that any two graphs G and H in X are isomorphic. By induction, H is isomorphic to every graph in the set  $X_1 = X \setminus G$  because  $|X_1| = k$ . Similarly, by induction, G is isomorphic to every graph in the set  $X_2 = X \setminus H$ . Let K be a graph in  $X_1 \cap X_2$ . So,  $G \cong K$  and  $H \cong K$ . Thus,  $G \cong H$  because  $\cong$  is transitive.

• The complement of a graph G = (V, E) is another graph  $G^{\circ} = (V, E^{\circ})$  with the same vertex set V and contains an edge for every  $\{u, v\}$  such that  $\{u, v\} \notin E$ . Let G be the cycle graph of length 5 with the vertex set  $\{0, 1, 2, 3, 4\}$ . The edges in G are the pairs  $\{\{i, j\} \mid j \equiv i + 1 \pmod{5}\}$ . Prove that  $G \cong G^{\circ}$  by giving the isomorphism.