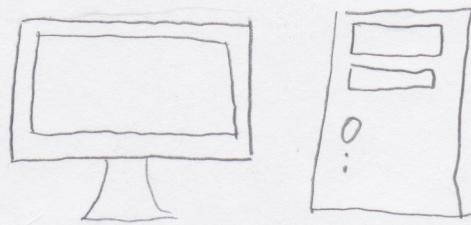
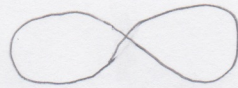


What is a Turing Machine?

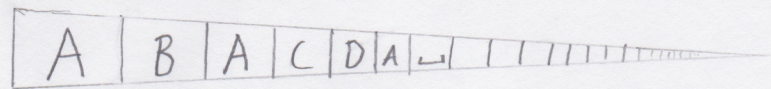
- It is a model of a general purpose computer.




- It has unlimited and unrestricted memory.




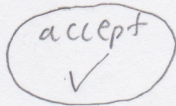
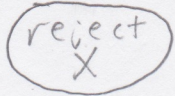
- This memory is stored in an infinitely long tape.

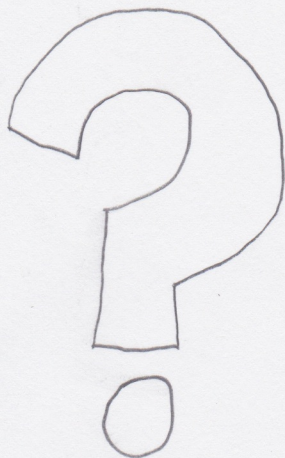


- A  points to the current position in memory of the Turing machine. The Turing machine can read or write at the position of the tape head.

- The tape head can move to the left or the right one symbol at a time.

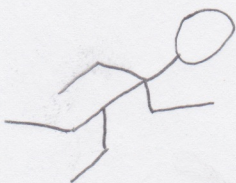
It must  moving to the left if the beginning of the tape is reached.

- A Turing machine will continue to run until either an  or  state is reached.



Why should we care?

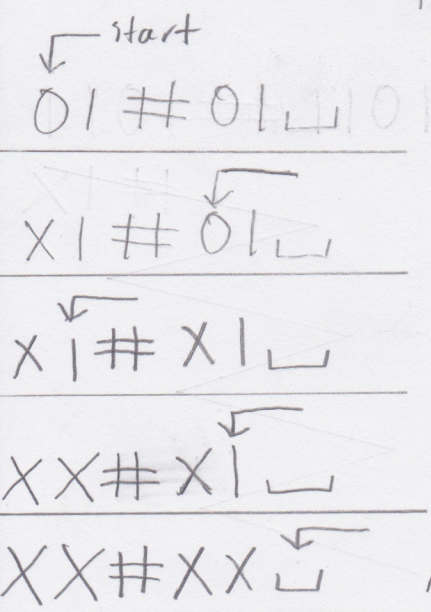
- A Turing machine can be created to solve any problem that is within the theoretical limits of computation.
- Turing machines are able to solve problems that cannot be solved using finite automata.



Running a Turing Machine

This example will show the strategy that a Turing machine would use to accept a language where all symbols before a # sign are the same as those after. Otherwise, the machine will reject.

Tape Configurations



Strategy

- start at the leftmost element, and mark it off with an X.
 - Move right until a # symbol is found.
 - Move to the right again. If this symbol is the same as the leftmost symbol was, mark it off as an X.
 - Now move to the left until the # sign is found. Then move left again until a symbol marked with an X is found. Mark off the symbol to the right of this.
 - Continue to Zig-Zag until either an unmatched pair is found, or all are Xed off.
- REJECT ACCEPT

What is a Turing Machine's Configuration?

- A setting of the following three items at a given point in time while the Turing machine is being run.

○ Current state

○ Current tape contents

A	A	B	␣
---	---	---	---

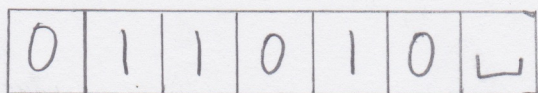
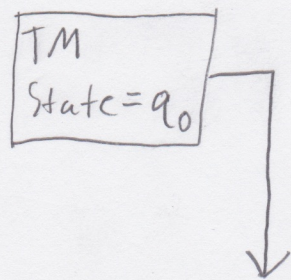
○ Current head location

▲

How can a configuration be represented?

- Using a string created in the following way:

Configuration \longrightarrow Representative String

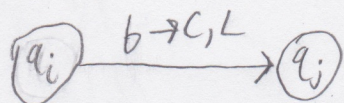


0110 q₀10
 ↑

The current state name is put to the left of the symbol the tape head is pointing at.

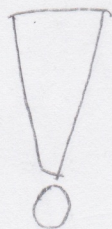
What is Yield?

- One configuration Yields another if the second configuration can be reached from the first in 1 step.
- Say the following transition exists in the state diagram representing TM M :



- Then we could say that configuration:

$aaq_i b v$ **YIELDS** $aaq_j a c v$



Special Configurations

- Start Configuration: $q_0 w$, where q_0 is the start state, and w is the initial contents of the tape.

★ - Accept Configuration: $\Sigma^* q_{\text{accept}} \Sigma^*$, which is any configuration consisting of valid symbols and the accept state.

★ - Reject Configuration: $\Sigma^* q_{\text{reject}} \Sigma^*$, which is any configuration consisting of valid symbols and the reject state.

These are HALTING configurations, which means that they do not yield any further configurations once reached.



How does a TM accept?

- To accept, a sequence of configurations must exist where:
 $C_1, C_2, \dots, C_{k-1}, C_k$ must exist where:

1. C_1 is the start configuration.
2. each configuration C_i yields C_{i+1}
3. C_k is an accepting configuration.

Turing Recognizable Vs. Turing Decidable

- First, define recognizable and decidable:

Recognizable: A Turing machine can accept, reject, or loop on input. A Turing machine recognizes a language if it accepts when a string in that language is inputted.

Decidable: A Turing machine decides a language if it recognizes that language and does not loop for any input.

Turing-Recognizable

A language is Turing-Recognizable if some Turing machine recognizes it.

Turing-Decidable

A language is Turing-Decidable if some Turing machine decides it. Notice that all Turing-Decidable languages are also Turing-Recognizable.